## OCR Computer Science A Level

### 2.3.1 Path Finding Algorithms <br> Concise Notes

## Specification:

- Dijkstra's shortest path algorithm
- $A^{*}$ algorithm


## Dijkstra's algorithm

- Dijkstra's algorithm finds the shortest path between two nodes in a weighted graph.
- Graphs are used as an abstraction for real life scenarios.
- Nodes and edges can represent different entities.
- Implemented using a priority queue, with the smallest distances being stored at the front of the list.
- Below is a step-by-step example of using Dijkstra's algorithm to find the shortest path between $A$ and $E$.


Priority Queue

$$
\begin{array}{|l|l|l|l|l|}
\hline \mathrm{C}=3 & \mathrm{D}=6 & \mathrm{~B}=7 & \mathrm{E}=\infty & \\
\hline
\end{array}
$$

## Step 1

- Starting from the root node (A), add the distances to all of the immediately neighbouring nodes ( $\mathrm{B}, \mathrm{C}, \mathrm{D}$ ) to the priority queue.
- The table below shoud be used to keep track of visited nodes and distances.
- Begin by filling in the 'Node' column with the nodes connected to the root node.
- The 'From' column should contain the node that you are travelling from.
- The 'Distance' column contains the distance between these two nodes.
- The 'Total distance' is the sum of the distances from the root node to that particular node.
- The node which is the shortest distance away from the root node is highlighted and selected for the next stage.

| Node | From | Distance | Total distance |
| :--- | :--- | :--- | :--- |
| B | A | 7 | 7 |
| C | A | 3 | 3 |
| D | A | 6 | 6 |


$A C=3$
Priority Queue

$$
\begin{array}{|l|l|l|l|l|}
\hline \mathrm{B}=5 & \mathrm{D}=6 & \mathrm{E}=13 & & \\
\hline
\end{array}
$$

## Step 2

- Remove the node the shortest distance away, from the front of the queue.
- Now traverse all of the nodes connected to the removed node C.
- If the total distance passing through the removed node to the neighbouring node is smaller than the distance currently stored with this node, update this value to the smaller distance.
- C has two neighbours: $B$ and $E$. The shortest total cost of travelling to $E$ so far is 13 , as it is less than the infinite value previously allocated to E .
- Travelling from $A$ to $C$ to $B$ adds up to a total cost of 5 , while travelling from $A$ to $B$ has a cost of 7 . B's cost value is therefore updated to 5 .
- We can ignore nodes we have visited, such as B.
- We can remove routes that are not the shortest way to get to a particular node.

| Node | From | Distance | Total distance |
| :--- | :--- | :--- | :--- |
| B | A | 7 | 7 |
| G | A | 3 | 3 |
| D | A | 6 | 6 |
| B | C | 2 | 5 |
| E | C | 10 | 13 |

Step 3

- Continue repeating Step 2 until the goal node has been reached.

| Node | From | Distance | Total distance |
| :--- | :--- | :--- | :--- |
| B | A | 7 | 7 |
| G | A | 3 | 3 |
| D | A | 6 | 6 |
| B | G | $Z$ | 5 |
| E | C | 10 | 13 |
| B | B | 3 | 8 |

- Once again, we repeat this same process for node D.

| Node | From | Distance | Total distance |
| :--- | :--- | :--- | :--- |
| B | A | 7 | 7 |
| G | A | 3 | 3 |
| D | A | 6 | 6 |
| B | G | $Z$ | 5 |
| E | G | 10 | 13 |
| B | B | 3 | 8 |
| E | D | 5 | 11 |

- As we have now visited all of the nodes on the graph, we can confirm by tracing back through the table above that the shortest path is ADE.

A* algorithm

- TA general path-finding algorithm which is an improvement of Dijkstra's algorithm and has two cost functions:

1) The actual cost between two nodes - the same cost as is measured in Dijkstra's algorithm.
2) An approximate cost from node $x$ to the final node, called a heuristic, which aims to make the shortest path finding process more efficient.

- The approximate cost is added onto the actual cost to determine which node is visited next.



## Step 1

- The method used here is very similar to the method used in Dijkstra's algorithm,
- The difference is that the heuristic cost is added onto the actual cost to calculate the total cost.
- Again, the route with the lowest total cost is selected to traverse further.

| Node | From | Distance | Heuristic | Total distance |
| :--- | :--- | :--- | :--- | :--- |
| B | A | 7 | 6 | 13 |
| C | A | 3 | 7 | 10 |
| D | A | 6 | 1 | 7 |

## Step 2

- The node D is then selected.
- Note that the heuristic cost of the previous node is not added on to the new distance.
- As 11 is the shortest total distance, the algorithm terminates.
- The shortest route is found to be ADE, at a cost of 11.

| Node | From | Distance | Heuristic | Total distance |
| :--- | :--- | :--- | :--- | :--- |
| B | A | 7 | 6 | 13 |
| C | A | 3 | 11 | 14 |
| D | A | 6 | 12 | 18 |
| E | D | 5 | - | 11 |

- The heuristics used here allow the shortest path to be found much quicker than when using Dijkstra's algorithm.
- How effective the $A^{*}$ algorithm is depends on the accuracy of the heuristics used.

